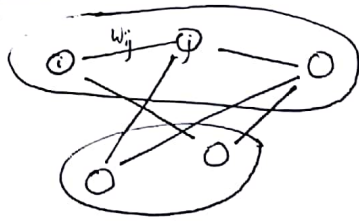


Core problem: mincut & variants



goal: minimize the weight of cut edges  
(subject to additional constraints)

Motivation: clustering of data

→ often, we get datapoints  $x_1, \dots, x_n$  with some notion of similarity:  $x_i$  &  $x_j$  similarity score is  $w_{ij}$ .

→ half of clustering is about MINIMIZING INTERCLUSTER WEIGHTS

• the other half is MAXIMIZING INTRACLUSTER WEIGHTS.

Note: many ways to construct this similarity matrix

- k-NN graph
- $\epsilon$ -neighborhood graph
- Gaussian kernel for Euclidean spaces

$$w_{ij} = \exp\left(-\frac{\|x_i - x_j\|^2}{2\sigma^2}\right)$$

Setting & Notation:

•  $G = (V, E)$  undirected graph

•  $W$  weights matrix

$w_{ij}$  is weight of an edge  $(i, j) \in E$

$w_{ij} = 0$  if  $(i, j) \notin E$

- Properties: -  $W$  is symmetric

- if  $w_{ij} = 1 \iff (i, j) \in E$ , then  $W$  is the usual adjacency matrix.

•  $D$  degree matrix =  $\text{diag}(d_1, \dots, d_n)$

$$d_i := \sum_{j \in V} w_{ij}$$

• if  $S \subset V$ ,  $\bar{S} := V \setminus S$

• if  $S, T \subset V$ ,  $E(S, T) := \{(i, j) \in E : i \in S, j \in T\}$

•  $\partial S = E(S, \bar{S})$

•  $\text{weight}(F) = \sum_{e \in F} w_e$  for  $F \subseteq E$

• Measure of size  $S \subseteq V$ :

-  $|S| := \text{cardinality}$

-  $\text{vol}(S) := \sum_{i \in S} d_i = \sum_{\substack{i \in S \\ j \in V}} w_{ij} = \text{weight}(E(S, V))$ .

MINCUT problem

$\min_{\emptyset \neq S \neq V} \text{weight}(E(S, \bar{S})) = \text{weight}(\partial S)$ .

Fact (Steiner-Wagner, 1995). There is an efficient algorithm that solves MINCUT.

QUESTIONS: is this good for clustering?

(No...)

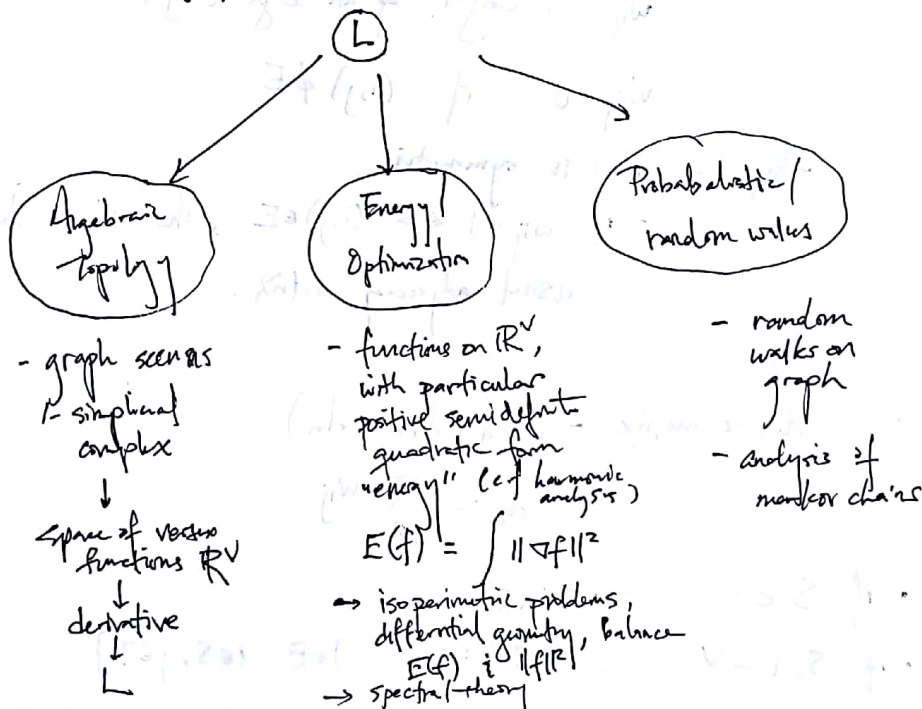
→ we can try to "balance" the size of each partition

→ turns out the problem becomes NP-hard (Wagner & Wagner 1993).

→ approach through spectral graph theory.

SPECTRAL GRAPH THEORY

At its heart, the graph Laplacian,  $L$ . Many connections



# REVIEW of LAPLACIAN

③

$$G = (V, E)$$

Dictionary between discrete calculus & vector calculus.

-  $V$  is space

$M$  is space

-  $\mathbb{R}^V$  is space of vertex functions

$C^\infty(M; \mathbb{R})$  space of functions

• inner product space

• inner product space

$$\langle f, g \rangle_V := \sum_{i \in V} f(i)g(i)$$

$$\langle f, g \rangle_M := \int_M f \cdot g$$

-  $\mathbb{R}^E$  is space of alternating vertex functions

$\mathcal{X}(M)$  space of vector fields

•  $X([i, j]) = -X([j, i])$  for  $(i, j) \in E$

• inner product space

$$\langle X, Y \rangle_E := \sum_{e \in E} w_e X(e) Y(e)$$

-  $d$  the derivative operator  $d: \mathbb{R}^V \rightarrow \mathbb{R}^E$

$\nabla$  gradient operator

$$df([i, j]) = f(j) - f(i)$$

- Dirichlet energy, measure of smoothness

$$E(f) := \langle df, df \rangle_E = \sum_{\substack{i, j \\ i \sim j}} w_{ij} (f(i) - f(j))^2$$

$$E(f) = \int_M \|\nabla f\|^2$$

Recall the adjoint:

if  $T: A \rightarrow B$  map of finite-dim inner-product spaces, then

$$\exists! T^*: B \rightarrow A \text{ s.t. } \forall a \in A, b \in B, \langle b, T a \rangle_B = \langle T^* b, a \rangle_A$$

- Laplacian  $L$  defined so that:

$$\langle df, df \rangle_E = \langle f, Lf \rangle_V$$

Why consider  $L$ ? Optimization problem:

$$\begin{aligned} \text{eg. } \min \langle df, df \rangle_E & \xrightarrow{\text{Lagrange multiplier}} Lf = \lambda f \\ \text{s.t. } \|f\|^2 &= 1 \end{aligned}$$

Properties of Laplacian:

①  $L = D - W$

③  $L \in \text{Sym}(\mathbb{R}^{n \times n})$

②  $\mathbf{1} \in \ker(L)$

④  $\forall A \subset V, \mathbf{1}_A^T L \mathbf{1}_A = \sum_{\substack{(i, j) \in E \\ i, j \in A}} w_{ij} (\mathbf{1}_A(i) - \mathbf{1}_A(j))^2 = \text{weight}(E(A, A))$

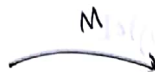
REVIEW of SPECTRAL THEORY: geometric & variational characterization of eigenvalues

Def. Let  $M \in \mathbb{R}^{n \times n}$ . An eigenvector  $v \in \mathbb{R}^n$  satisfies  $v \neq 0, \exists \lambda \in \mathbb{R}$  s.t.  
 $Mv = \lambda v$ .

The eigenvalues of a matrix form its spectrum.

Thm (Spectral Theorem for Sym. Mat).

Geometric characterization:



Def. Let  $M \in \mathbb{R}^{n \times n}$ . Let  $v \in \mathbb{R}^n - \{0\}$ . Rayleigh quotient

$$R(v) = \frac{v^T M v}{v^T v}$$

notice:  $R(\alpha v) = R(v) \quad \forall \alpha \in \mathbb{R} - \{0\}$ , can assume  $v \in S^{n-1}$ .

Thm (Variational Characterization). Let  $M \in \text{Sym}(\mathbb{R}^{n \times n})$

$\lambda_1 \leq \dots \leq \lambda_n$  eigenvalues of  $M$ .

$$\lambda_1 = \min_{v \in S^{n-1}} R(v) \quad , \text{ let } v_1 = \text{argmin}$$

!

$$\lambda_i = \min_{v \in S^{n-1}} R(v)$$

$$v \perp \{v_1, \dots, v_{i-1}\}$$

Pf. minimize  $v^T M v$  s.t.  $\|v\|^2 = 1$ .

$$\Rightarrow Mv = \lambda v, \text{ and } v^T M v = \lambda v^T v = \lambda.$$

Exercise Let  $M \in \text{Sym}(\mathbb{R}^{n \times n})$ . Let  $X = \begin{bmatrix} | & & | \\ x_1 & \dots & x_k \\ | & & | \end{bmatrix} \quad x_i \in \mathbb{R}^n$

$$\min_X \sum_{i=1}^k x_i^T M x_i \quad \text{s.t. } x_i \perp x_j \text{ when } i \neq j$$

$$\|x_i\|^2 = 1$$

$$\Leftrightarrow \min_X \text{Tr}(X^T M X) \quad \text{s.t. } X^T X = I_{k \times k}$$

QUESTION: is  $X$  unique?

No. Let  $U$  be unitary matrix in  $\mathbb{R}^{k \times k}$ .

⑤

Fact.  $\text{Tr}(ABC) = \text{Tr}(BCA)$ .

$$\begin{aligned}\text{Tr}((XU)^T M (XU)) &= \text{Tr}(M (XU)(XU)^T) \\ &= \text{Tr}(M X U U^T X^T) \\ &= \text{Tr}(M X X^T) = \text{Tr}(X^T M X).\end{aligned}$$

$$(XU)^T (XU) = U^T X^T X U = U^T U = \text{Id}_{k \times k}.$$

$\Rightarrow$  Equivalently,

$$\min_X \text{Tr} \left( X^T \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_k \end{bmatrix} X \right) \quad (\lambda_1 \leq \dots \leq \lambda_k).$$

s.t.  $X^T X = \text{Id}_{k \times k}$

$\rightarrow$  minimized at  $\sum_{i=1}^k \lambda_i$

$$\Rightarrow X^* = \begin{bmatrix} | & & | \\ v_1 & \dots & v_k \\ | & & | \end{bmatrix}$$

$v_i$  correspondingly eigenvectors

Exercise (Generalized Eigenvector Problem). Let  $M, N \in \text{Sym}(\mathbb{R}^{n \times n})$ ,  $N \geq 0$ .

What is solution to eigenvalue problem?

$$Mv = \lambda Nv$$



$$N^{-1/2} M N^{1/2} N^{1/2} v = \lambda N^{1/2} v$$

Thus, equivalent to solving the eigenvalue problem for  $N^{-1/2} M N^{-1/2}$ .

### Preview of Laplacians

We saw that  $L = D - W$ .

- hint: we want to solve optimization problems ~~that~~ <sup>with</sup> constraints

$$Lv = \lambda Dv$$

$\Rightarrow$  define  $D^{-1}L =: L_{rw}$  the random walk Laplacian,  $L_{rw} = I - D^{-1}W$

- hint: eigenvalues will be equivalent to solving eigenvalue problems for

$$\begin{aligned}D^{1/2} L D^{-1/2} &=: L_{sym} \quad \text{the } \underline{\text{symmetrized/normalized Laplacian}} \\ &= I - D^{1/2} W D^{-1/2}\end{aligned}$$

## RETURN TO GRAPH CUT : the probabilistic version

- give graph a random walk interpretation.

- a random walker transitions from a vertex  $i$  to adjacent vertex  $j$  w.p.

$$\frac{w_{ij}}{\sum_k w_{ik}} = \frac{w_{ij}}{d_i}$$

Exercise let  $\pi^{(t)} = [\pi_1, \dots, \pi_n]$  probability mass of a random walker at time  $t$ .

After 1 time step, what is the new  $\pi^{(t+1)}$ ?

$$\pi_j^{(t+1)} = \sum_i \pi_i \cdot \frac{w_{ij}}{d_i} = (\pi D^{-1} W)_j$$

Thus,  $P = D^{-1}W$  is the transition matrix.

## Spectral Clustering

Goal: Cut  $V$  into two sets  $A \sqcup \bar{A}$ . Let random walker on graph. Can we minimize

$$\Pr[A \rightarrow \bar{A} \text{ or } \bar{A} \rightarrow A]$$

the probability the random walker crosses from  $A$  into  $\bar{A}$  or vice versa?

Exercise: Compute  $\Pr[A \rightarrow \bar{A}]$

$$= \frac{\text{weight}(E(A, \bar{A}))}{\text{weight}(E(A, V))} = \frac{\mathbb{1}_A^T L \mathbb{1}_{\bar{A}}}{\sum_{i \in A} d_i} = \frac{\mathbb{1}_A^T L \mathbb{1}_A}{\text{vol}(A)}$$

$$\text{Thus, } \Pr[A \rightarrow \bar{A} \text{ or } \bar{A} \rightarrow A] = \frac{E(A, \bar{A})}{\text{vol}(A)} + \frac{E(\bar{A}, A)}{\text{vol}(A)}$$

$$\text{Define } N_{\text{cut}}(A_1, \dots, A_k) = \sum_{i=1}^k \frac{E(A_i, \bar{A}_i)}{\text{vol}(A_i)}$$

$$V = A_1 \sqcup \dots \sqcup A_k$$

Problem:  $\min_A N_{\text{cut}}(A, \bar{A}), \quad 0 \neq A \neq V$

Fact. This is NP-hard.

# Optimization Problem

$$\min_A \underbrace{\frac{\mathbb{1}_A^T L \mathbb{1}_A}{\text{vol}(A)} + \frac{\mathbb{1}_{\bar{A}}^T L \mathbb{1}_{\bar{A}}}{\text{vol}(\bar{A})}}$$

~~$$\frac{\mathbb{1}_A^T}{\sqrt{\text{vol}(A)}} L \frac{\mathbb{1}_A}{\sqrt{\text{vol}(A)}}$$~~

$$f_A := \frac{\mathbb{1}_A^T}{\sqrt{\text{vol}(A)}}$$

Observe:

$$A \sqcup \bar{A} \Leftrightarrow \mathbb{1}_A + \mathbb{1}_{\bar{A}}$$

$$f_A^T D f_A = \sum_{i \in A} d_i \left( \frac{1}{\sqrt{\text{vol}(A)}} \right)^2 = \frac{1}{\text{vol}(A)} \sum_{i \in A} d_i = 1.$$

$$\Rightarrow \begin{bmatrix} f_A & f_{\bar{A}} \end{bmatrix}^T D \begin{bmatrix} f_A & f_{\bar{A}} \end{bmatrix} = I_2.$$

Let  $X = \begin{bmatrix} f_A & f_{\bar{A}} \end{bmatrix}$

⇒ Optimize:

$$\min_{X \in \{ \begin{bmatrix} f_A & f_{\bar{A}} \end{bmatrix} : A \subset V \}} \text{Tr}(X^T L X) \quad \text{s.t. } X^T D X = I_2$$

↓ Relax

$$\min_X \text{Tr}(X^T L X) \quad \text{s.t. } X^T D X = I_2.$$

We saw that the solution corresponds to

$$Lx = \lambda D x$$

↑

$$L_{\text{sym}}(D^{1/2} x) = \lambda (D^{1/2} x)$$

⇒ if  $v$  is eigenvector of  $L_{\text{sym}}$ , then  $D^{-1/2} v$  is solution to generalized eigenvalue problem.

$$\Rightarrow X = \begin{bmatrix} D^{-1/2} v_1 & D^{-1/2} v_2 \end{bmatrix} \quad (\text{note: } D^{-1/2} v_i = \mathbb{1} \text{ for } L_{\text{sym}}).$$

# Graph Clustering (Shi & Malik 2000) Ncut.

Input  $n$  data points  $x_1, \dots, x_n$   
 $W$  similarity matrix

Output  $A_1 \cup \dots \cup A_k = V$  partition of  $x_i$ 's

Goal  $\min \text{Tr}(X^T L X)$  s.t.  $X^T D X = I$   
 $X \in \{0, 1\}^{n \times k}$

Relax  $\min_{X \in \mathbb{R}^{n \times k}} \text{Tr}(X^T L X)$  s.t.  $X^T D X = I$

$$\Leftrightarrow Lx_i = \lambda_i D x_i \Leftrightarrow x_i = D^{-1/2} v_i$$

$v_i \in \text{Spectrum}(L_{\text{sym}})$ .

## Algorithm.

1. Compute the  $k$  generalized eigenvectors  $x_1, \dots, x_k$   
 $Lx_i = \lambda_i D x_i$ .

2. Represent  $x_i$  by  $i$ th row of  $X$

$$X = \begin{bmatrix} | & & | \\ x_1 & \dots & x_k \\ | & & | \end{bmatrix} \leftarrow x_i \text{ represented in } \mathbb{R}^k$$

3. Run  $k$ -means on the rows of  $X$ .

Note. HW1's spring problem is equivalent to the RatioCut problem.

$$\text{RatioCut: } \min_{A_1 \cup \dots \cup A_k} \sum \frac{\mathbb{1}_{A_i}^T L \mathbb{1}_{A_i}}{|A_i|} \leftarrow \text{instead of } \text{vol}(A_i)$$

Prop. (Quatterly & Miller 1998)  $\forall C > 0, \exists G$  s.t.

$$C \cdot \text{OPT}(G) \leq \text{OPT}_{\text{approx}}(G)$$

$$\text{where } \text{OPT}(G) = \min_{\emptyset \neq A \subseteq V} \text{RatioCut}(A, \bar{A})$$

$$\text{OPT}_{\text{approx}} = \min \text{Convex Relaxation}$$

"Cockroach graph" constructive example.

Prop. (Bui & Jones 1992). Approx. problem NP-hard.